

Lecture 4. Lower bounds on Ramsey numbers

4.1 Lovász Local Lemma

Theorem 4.1 (Lovász Local Lemma (Symmetric Version)). *Let $\{A_i\}_{i \in [k]}$ be a family of random events. For any i , $\mathbb{P}[A_i] \leq p$. Any event is independent of all other events except for d of them. If $ep(d-1) < 1$, then the probability that all events' complements occur simultaneously is greater than 0, that is:*

$$\mathbb{P} \left[\bigcap A_i^c \right] > 0.$$

We focus on the asymmetric version of the Lovász Local Lemma, which is stronger than the symmetric version. Firstly, we define the following auxiliary graph.

Definition 4.2. Let $\mathcal{A} = \{A_1, \dots, A_k\}$ be a family of random events in the probability space Ω . Let $D := D_{\mathcal{A}}$ be the *dependence graph* with vertex set $V(D) = \mathcal{A}$ and edge set $E(D) = \{A_i A_j : A_i \text{ and } A_j \text{ are dependent for each } i, j \in [k]\}$.

Theorem 4.3 (Lovász Local Lemma (Asymmetric Version)). *Given a probability space (Ω, \mathbb{P}) , the event collection \mathcal{A} and the dependence graph D in Definition 4.2. Denote the neighborhood set of A_i in D by $N(A_i)$. If there exists a mapping $f : \mathcal{A} \rightarrow [0, 1)$ satisfying that*

$$\mathbb{P}[A_i] \leq f(A_i) \prod_{B \in N(A_i)} (1 - f(B))$$

holds for each $A_i \in \mathcal{A}$, then the following holds:

$$\mathbb{P} \left[\bigcap A_i^c \right] > 0.$$

4.2 Applications of Lovász Local Lemma

The Ramsey number $r(k, \ell)$ is the smallest integer N such that any red-blue edge-coloring of K_N contains a red K_k or a blue K_ℓ .

Remark:

- Ramsey's Theorem [3]: The Ramsey number exists.
- Erdős and Szekeres [2]: $r(k, \ell) \leq \binom{k+\ell-2}{k-1}$. If $k = \ell$, then this yields $r(k, k) \leq 4^k$.
- Campos, Griffiths, Morris and Sahasrabudhe [1]: There exists $\varepsilon > 0$ such that $r(k, k) \leq (4 - \varepsilon)^k$.

Theorem 4.4.

$$r(3, \ell) = \Omega \left(\frac{\ell^2}{\log^2 \ell} \right).$$

Proof. Consider a random edge-coloring of the complete graph K_n , where each edge is colored red with probability p and blue with probability $1 - p$ independently. Our goal is to obtain that with positive probability there is a coloring without a red triangle and without a blue K_ℓ , since this would establish the lower bound $r(3, \ell) > n$.

For each 3-element set $T \subseteq [n]$, let A_T be the event that T induces a red K_3 . Note that for each T , we have $\mathbb{P}[A_T] = p^3$. For each ℓ -element set $S \subseteq [n]$, let B_S be the event that S induces

a blue K_ℓ . Note that for each ℓ , we have $\mathbb{P}[B_S] = (1-p)^{\binom{\ell}{2}}$. Let us now define a dependence graph D for these events. We join two events of the form A_T or B_S , if the corresponding sets S or T share an edge. Now we can bound the degrees in this graph,

$$\begin{cases} \deg(A_T) \leq \binom{3}{2}(n-3) + \binom{3}{2}\binom{n-3}{\ell-2} \leq 3n + \binom{n}{\ell} \text{ for each } A_T, \\ \deg(B_S) \leq \binom{\ell}{2}n + \binom{\ell}{2}\binom{n-3}{\ell-2} \leq \binom{\ell}{2}n + \binom{n}{\ell} \text{ for each } B_S. \end{cases}$$

Define a mapping $f : \{A_T : T \in \binom{[n]}{3}\} \cup \{B_S : S \in \binom{[n]}{\ell}\} \rightarrow [0, 1)$ such that $x := f(A_T)$ and $y := f(B_S)$. In order to use Lovász Local Lemma, we need to find positive real numbers $x, y \in (0, 1)$ such that

$$\begin{cases} p^3 \leq x(1-x)^{3n}(1-y)^{\binom{n}{\ell}}, \\ (1-p)^{\binom{\ell}{2}} \leq y(1-x)^{n\binom{\ell}{2}}(1-y)^{\binom{n}{\ell}}, \end{cases}$$

Let us now try to find such $p, x, y \in (0, 1)$ for sufficiently large n . We choose $y = \frac{1}{\binom{n}{\ell}}$, then $(1-y)^{\binom{n}{\ell}} \approx \frac{1}{e}$. Furthermore, we observe that p and x need to fulfill the following inequalities:

$$p^3 \leq x(1-x)^{3n}(1-y)^{n/2} \leq x,$$

$$e^{-p\binom{\ell}{2}} \approx (1-p)^{\binom{\ell}{2}} \leq y(1-x)^{\binom{\ell}{2}n}(1-y)^{n/2} \leq (1-x)^{\binom{\ell}{2}n} \approx e^{-xn\binom{\ell}{2}}.$$

Hence we need $p \geq xn \geq p^3n$. Therefore $p \leq \frac{1}{\sqrt{n}}$ and $x \geq p^3$. Finally, for the second condition, we note

$$e^{-p\binom{\ell}{2}} \approx (1-p)^{\binom{\ell}{2}} \leq y(1-x)^{\binom{\ell}{2}n}(1-y)^{n/2} \leq y = \frac{1}{\binom{n}{\ell}} \approx e^{-\ell \log n},$$

hence $p\ell^2 \geq p\binom{\ell}{2} \geq \ell \log n$ and therefore $\ell \geq \frac{1}{p} \log n \geq \sqrt{n} \log n$.

Motivated by this we may assume $\ell \geq 20\sqrt{n} \log n$ and choose $y = \frac{1}{\binom{n}{\ell}}$, $x = \frac{1}{9n^{3/2}}$ and $p = \frac{1}{3\sqrt{n}}$. After choosing the constants, we can give the full proof. For sufficiently large n , we have

$$(1-y)^{\binom{n}{\ell}} = \left(1 - \frac{1}{\binom{n}{\ell}}\right)^{\binom{n}{\ell}} \geq e^{-1.01},$$

$$(1-x)^{3n} = \left(1 - \frac{1}{9n^{3/2}}\right)^{3n} \geq 1 - \frac{1}{3\sqrt{n}} \geq e^{-0.01}.$$

Thus,

$$p^3 = \frac{1}{27n^{3/2}} \leq \frac{1}{9n^{3/2}} \cdot \frac{1}{3} \leq \frac{1}{9n^{3/2}} e^{-1.02} \leq x(1-x)^{3n}(1-y)^{\binom{n}{\ell}},$$

which establishes the first desired inequality.

For the second inequality, for sufficiently large n , we get

$$(1-x)^{\binom{\ell}{2}n} \geq e^{-2xn\binom{\ell}{2}} \geq e^{-\frac{2}{9\sqrt{n}}\binom{\ell}{2}}.$$

Furthermore, using $\ell \geq 20\sqrt{n} \log n$, we have

$$y = \frac{1}{\binom{n}{\ell}} \geq \frac{1}{n^\ell} = e^{-\ell \log n} \geq e^{-\ell^2 \frac{1}{20\sqrt{n}}} \geq e^{-\ell(\ell-1) \frac{1}{19\sqrt{n}}} \geq e^{-\frac{1}{9\sqrt{n}}\binom{\ell}{2} + 1.01}.$$

Hence

$$(1-p)^{\binom{\ell}{2}} \leq e^{-p\binom{\ell}{2}} = e^{-\frac{1}{3\sqrt{n}}\binom{\ell}{2}} = e^{-\frac{1}{9\sqrt{n}}\binom{\ell}{2} + 1.01} e^{-\frac{2}{9\sqrt{n}}\binom{\ell}{2}} e^{-1.01} \leq y(1-x)^{\binom{\ell}{2}n}(1-y)^{\binom{n}{\ell}},$$

which verifies the second desired inequality.

By Lemma 4.3, we obtain that

$$\mathbb{P} \left[\bigcap_{T \in \binom{[n]}{3}} A_T^c \cap \bigcap_{S \in \binom{[n]}{\ell}} B_S^c \right] > 0.$$

So for $\ell \geq 20\sqrt{n} \log n$, we can find p, x, y such that there exists a 2 edge-coloring of K_n such that there is no red triangle and there is no blue K_ℓ . This implies $r(3, \ell) > n$.

Note that $n \leq \frac{\ell^2}{(40 \log \ell)^2}$ implies

$$20\sqrt{n} \log n \leq 20 \frac{\ell}{40 \log \ell} \log \ell^2 = \ell.$$

Therefore we have $r(3, \ell) \geq \frac{\ell^2}{(40 \log \ell)^2}$. ■

Theorem 4.5 (Erdős' Lower Bound). *Let $C \geq 1$ and $p_C \in (0, 1/2]$ be the unique solution to $C = \frac{\log p_C}{\log(1-p_C)}$. Let $M_C = p_C^{-1/2}$. Then $r(\ell, C\ell) = \Omega(\ell \cdot M_C^\ell)$. In particular, when $C = 1$, we have $p_C = 1$ and $r(\ell, \ell) = \Omega(\ell \sqrt{2}^\ell)$.*

Proof. Let $p \in (0, 1/2]$. Consider a random edge-coloring of the complete graph K_n , where each edge is independently colored red with probability p and blue with probability $1 - p$. Let

$$f(n, p) := A(n, p) + B(n, p) \text{ where } A(n, p) = \binom{n}{\ell} p^{\binom{\ell}{2}} \text{ and } B(n, p) = \binom{n}{C\ell} (1-p)^{\binom{C\ell}{2}}.$$

Note that

$$\mathbb{P}[\text{There exists a red } K_\ell \text{ or a blue } K_{C\ell}] \leq f(n, p).$$

Hence, if $f(n, p) = 1 - o_\ell(1)$, then there exists at least one such coloring with no red K_ℓ and no blue $K_{C\ell}$, implying $r(\ell, C\ell) > n$. It thus suffices to find the maximum value of $n = n(p)$ such that $f(n, p) = o_\ell(1)$. Assume this maximum is achieved at $p_0 = p_{C, \ell}$. Then,

$$\frac{\partial f(n, p_0)}{\partial p} = \frac{\binom{\ell}{2}}{p_0} \binom{n}{\ell} p_0^{\binom{\ell}{2}-1} - \frac{\binom{C\ell}{2}}{1-p_0} \binom{n}{C\ell} (1-p_0)^{\binom{C\ell}{2}-1} = 0.$$

Thus, we have $\log A(n, p_0) = \log B(n, p_0) + O(\log \ell)$. Solving this along with $A(n, p_0) + B(n, p_0) = 1 - o_\ell(1)$, we obtain that $\log A(n, p_0) = O(\log \ell)$ and $\log B(n, p_0) = O(\log \ell)$. Therefore, we obtain that

$$\begin{cases} -\log p_0 = \frac{2 \log(en/\ell)}{\ell-1} + O\left(\frac{\log \ell}{\ell^2}\right), \\ -\log(1-p_0) = \frac{2 \log(en/C\ell)}{C\ell-1} + O\left(\frac{\log \ell}{\ell^2}\right). \end{cases}$$

We then derive that $p_0 = p_C + O(1/\ell)$, where the constant p_C satisfies $C = \frac{\log p_C}{\log(1-p_C)}$. It follows directly from the above that $n = \frac{\ell}{e} \cdot p_0^{-(\ell-1)/2} \cdot e^{O(\frac{\log \ell}{\ell})} = \Theta(\ell) \cdot (p_C + O(1/\ell))^{-\ell/2} = \Theta(\ell \cdot M_C^\ell)$, where $M_C := p_C^{-1/2}$. This establishes $r(\ell, C\ell) = \Omega(\ell \cdot M_C^\ell)$. ■

References

- [1] M. Campos, S. Griffiths, R. Morris, and J. Sahasrabudhe. An exponential improvement for diagonal Ramsey. *arXiv:2303.09521*, 2025.
- [2] P. Erdős and G. Szekeres. A combinatorial problem in geometry. *Compos. Math.*, 2:463–470, 1935.
- [3] F. P. Ramsey. On a problem of formal logic. *Proc. London Math. Soc.*, 30:264–286, 1930.

Extensive reading

The following two papers provide relevant extensive reading materials:

1. Noga Alon, Lajos Rónyai and Tibor Szabó, Norm-graphs: variations and applications, *Journal of Combinatorial Theory, Series B* 76 (1999), 280–290.
2. B. Bukh and D. Conlon, Rational exponents in extremal graph theory, *Journal of the European Mathematical Society*, 20 (2018), 1747–1757.